

Figure 1: Euler's problem of the bridges and a suitable graph

## **1** Principles of Graph Theory

With his problem of the bridges of Königsberg, Leonhard Euler (1707 - 1783) established the mathematical discipline of graph theory. The elements of a graph are vertices and edges. The graph on the right in figure 1 shows four vertices which represent the four regions of mainland and seven edges which represent the seven bridges over the river Pregel.

**Definition A.1** A *directed graph* or *digraph*  $\mathbf{G}$  is an ordered pair  $(\mathbf{V}, \mathbf{E})$ , where  $\mathbf{V}$  is a (finite or denumerably infinite) set of vertices and  $\mathbf{E}$  is a set of edges.

The *incidence function*:  $\mathbf{E} \to (\mathbf{V} \times \mathbf{V})$  assigns every edge  $\mathbf{e} \in \mathbf{E}$  an ordered pair  $(\mathbf{u}, \mathbf{v}) \in (\mathbf{V} \times \mathbf{V})$ . The pair  $\mathbf{e} \equiv (\mathbf{u}, \mathbf{v})$  is called a *directed edge*. We say  $\mathbf{e}$  and  $\mathbf{u}$  are *incident* and  $\mathbf{e}$  and  $\mathbf{v}$  are *incident*. We call  $\mathbf{G}$  finite, if  $\mathbf{V}$  and  $\mathbf{E}$  are finite sets.

The two endpoints of an edge  $(\mathbf{u}, \mathbf{v})$  are called *initial endpoint* (tail)  $\mathbf{p}_{\text{init}} := \mathbf{u}$ and *terminal endpoint* (head)  $\mathbf{p}_{\text{term}} := \mathbf{v}$ . In a diagram an arrowhead is placed at  $\mathbf{p}_{\text{term}}$ . Two edges are called *adjacent* if they have an endpoint in common. Two vertices are called *adjacent* ( $\sim$ ), if there is an edge which is incident to both. An edge with identical endpoints is called *loop*.

**Definition A.2** If every edge of a graph **G** is undirected, i. e. of the form  $\{\mathbf{u}, \mathbf{v}\}$ , then **G** is called an *undirected graph*.

**Definition A.3** A graph is called *simple*, if for any  $\mathbf{u}, \mathbf{v} \in \mathbf{V}$  we have

(S1) at most one of the pairs  $\mathbf{e} := (\mathbf{u}, \mathbf{v}), \ \bar{\mathbf{e}} := (\mathbf{v}, \mathbf{u}) \text{ or } \mathbf{f} := \{\mathbf{v}, \mathbf{u}\}$  is an element of  $\mathbf{E}$  (no *multiple edges*) and

(S2) a pair of the form  $(\mathbf{v}, \mathbf{v})$  or  $\{\mathbf{v}, \mathbf{v}\}$  is not an element of **E** (no *loops*).

**Definition A.4** An *oriented graph* is a simple directed graph.

**Definition A.5** A simple undirected graph is an oriented graph, the orientations of whose edges are abolished, i. e. every edge is a set of the form  $\{\mathbf{u}, \mathbf{v}\}$ .

On the other hand we get an oriented graph by assigning every edge of an unoriented simple graph a direction. However this is ambiguous. It can be done in  $2^{|\mathbf{E}|}$  different ways. Thus there are  $2^{|\mathbf{E}|}$  different oriented graphs over the same simple unoriented graph.

**Degrees of a vertex A.6** The *degree*  $deg(\mathbf{v})$  of a vertex  $\mathbf{v}$  is the number of edges that are incident to  $\mathbf{v}$ . Loops are counted twice. A vertex of degree 0 is called a *lone vertex*. For every vertex  $\mathbf{v}$  of a directed graph we distinguish between its *exit degree*  $deg^+(\mathbf{u}) = |\{\mathbf{v}|(\mathbf{u}, \mathbf{v}) \in \mathbf{E}\}|$  and its *initial degree*  $deg^-(\mathbf{u}) = |\{\mathbf{v}|(\mathbf{v}, \mathbf{u}) \in \mathbf{E}\}|$ .

We have that  $d^+(\mathbf{u})$  equals the number of initial vertices in a given graph. Since every directed edge has exactly one initial vertex, this number equals  $|\mathbf{E}|$ . Similarly  $d^-(\mathbf{u}) = |\mathbf{E}|$ , since every directed edge has exactly one terminal vertex:

$$\sum_{\mathbf{u}\in\mathbf{V}}\deg^+(\mathbf{u})=\sum_{\mathbf{u}\in\mathbf{V}}\deg^-(\mathbf{u})=|\mathbf{E}|.$$

Lemma A.7 (Handshaking Lemma) In any directed graph G = (V, E):

$$\sum_{\mathbf{u}\in\mathbf{V}}\deg(\mathbf{u})=2|\mathbf{E}|.$$

Corollary A.8 Any directed graph  $\mathbf{G} = (\mathbf{V}, \mathbf{E})$  has an even number of vertices which are of odd degree  $\deg(\mathbf{u})$ .

